

Preliminary Notes on FFT (helpful
in the solution of
Prob SP-1)

Let $x[n]$ be a N -point finite sequence
with $N = 2^P$, i.e. (1)

$$x[n] = \{x[0], x[1], x[2], \dots, x[N-2], x[N-1]\} \quad (2)$$

Then $x[n]$ can be broken up
into two sequences, called
even subsequence $x_1[n]$ and
odd subsequence $x_2[n]$

where

$$\left\{ \begin{array}{l} x_1[0] = x[0] \\ x_1[1] = x[2] \\ \vdots \\ x_1[m] = x[2m] \\ \vdots \\ x_1[\frac{N}{2}-1] = x[N-2] \end{array} \right\} \quad (3)$$

$$\left\{ \begin{array}{l} x_2[0] = x[1] \\ x_2[1] = x[3] \\ \vdots \\ x_2[m] = x[2m+1] \\ \vdots \\ x_2[\frac{N}{2}-1] = x[N-1] \end{array} \right\} \quad (4)$$

Solutions
to Hmk # 8 (cont)

Now let $x_1[n]$ and $x_2[n]$ be two N -pt. signals

$X[k]$ denote the N -pt. DFT of $x[n]$ (5)

$X_1[k]$ " " $\frac{N}{2}$ -pt DFT of $x_1[n]$ (6)

$X_2[k]$ " " " " " $x_2[n]$ (7)

Then the fundamental computational benefit of the FFT algorithm is that it allows one to obtain the components of $X[k]$ from those of $X_1[k]$ and $X_2[k]$ by the simple recursion

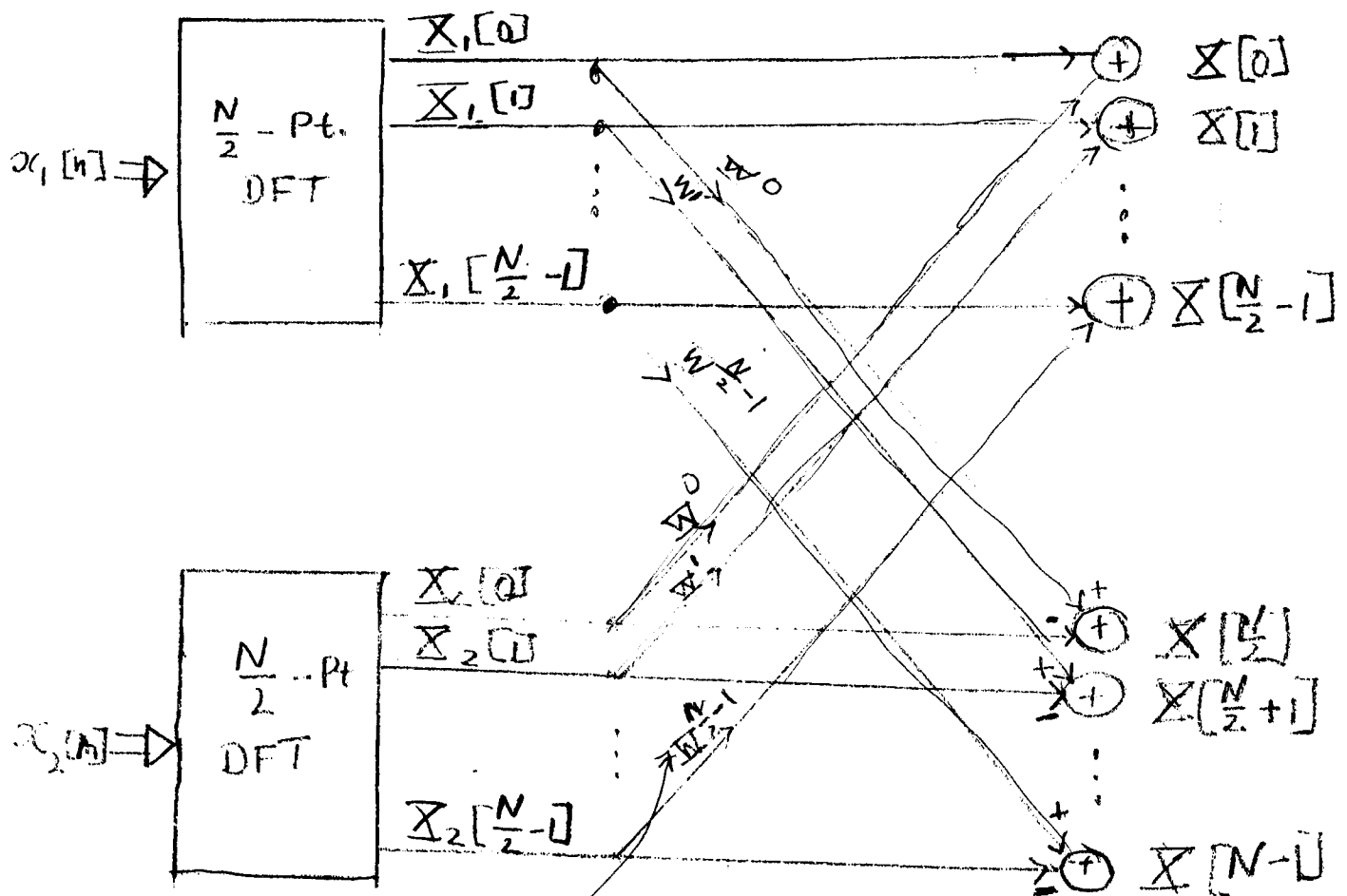
$$X[k] = \begin{cases} X_1[k] + W_N^k X_2[k], & k = 0, 1, \dots, \frac{N}{2} - 1 \quad (8a) \\ X_1[k - \frac{N}{2}] - W_N^{k - \frac{N}{2}} X_2[k - \frac{N}{2}], & k = \frac{N}{2}, \dots, N-1 \quad (8b) \end{cases}$$

Note that in 8(b) I used the fact that $W_N^k = -W_N^{k - \frac{N}{2}}$ (9)

which allows us to use only $\frac{N}{2}$ weights in this iteration

This procedure is exhibited in the below block diagram/signal flow graph. For simplicity, we abbreviate

$$\overline{W}_N = W \quad (9)$$



$W_N^{\frac{N}{2}-1}$

Fig 1

(Note that in the lower components, $k = \frac{N}{2}, \dots, N-1$, you subtract the weights)

Soln. to Hmk 8 (Cont)

Another way of representing Fig 1 would be through the symbolic representation in Fig 2 done in class and in Notes #6

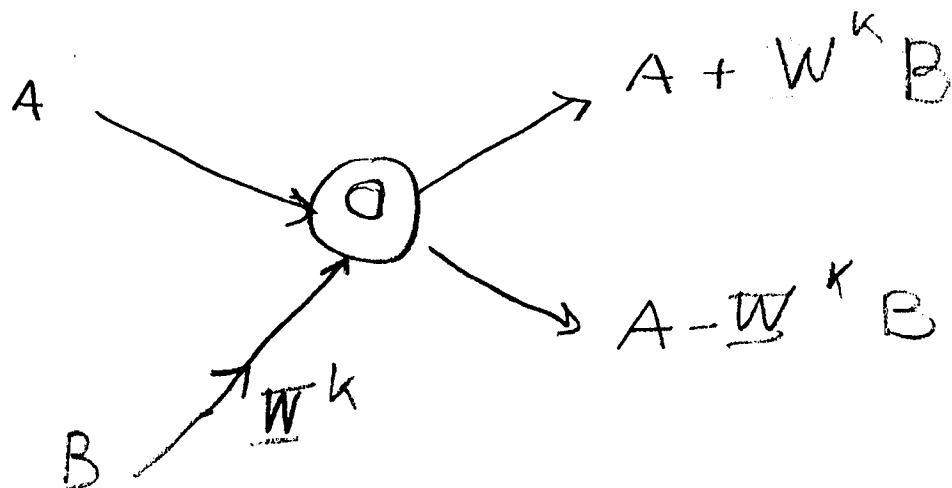


Fig. 2

This representation allows to represent Fig. 1 as Fig 3 on the next page

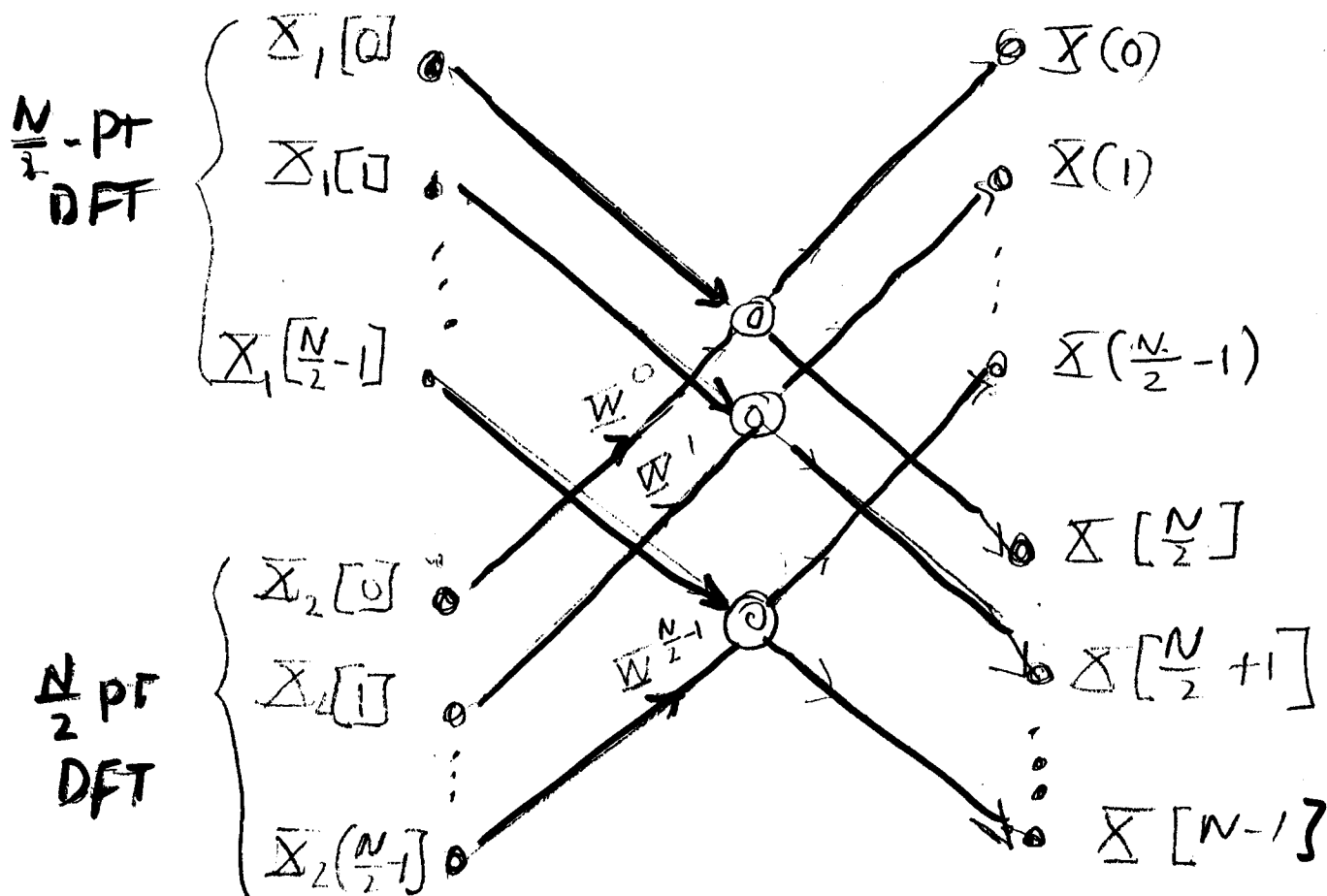


Fig 3

For an eight point FFT
see The Figure in p. 5
of Lec. Notes # 6.

NOTE

To easily obtain the arrangement
of data for 2 PT DFT, use
bit-reversal as indicated on the next page.

Solns. to Hunk 8 (Cont)

Hunk 8

6

Re-Ordering of data by

bit reversal	Forward order	Reverse Order	n
0	0 0 0	0 0 0	0 } x[0]
1	0 0 1	1 0 0	4 } x[4]
2	0 1 0	0 1 0	2 } x[2]
3	0 1 1	1 1 0	6 } x[6]
4	1 0 0	0 0 1	1 } x[1]
5	1 0 1	1 0 1	5 } x[5]
6	1 1 0	0 1 1	3 } x[3]
7	1 1 1	1 1 1	7 } x[7]

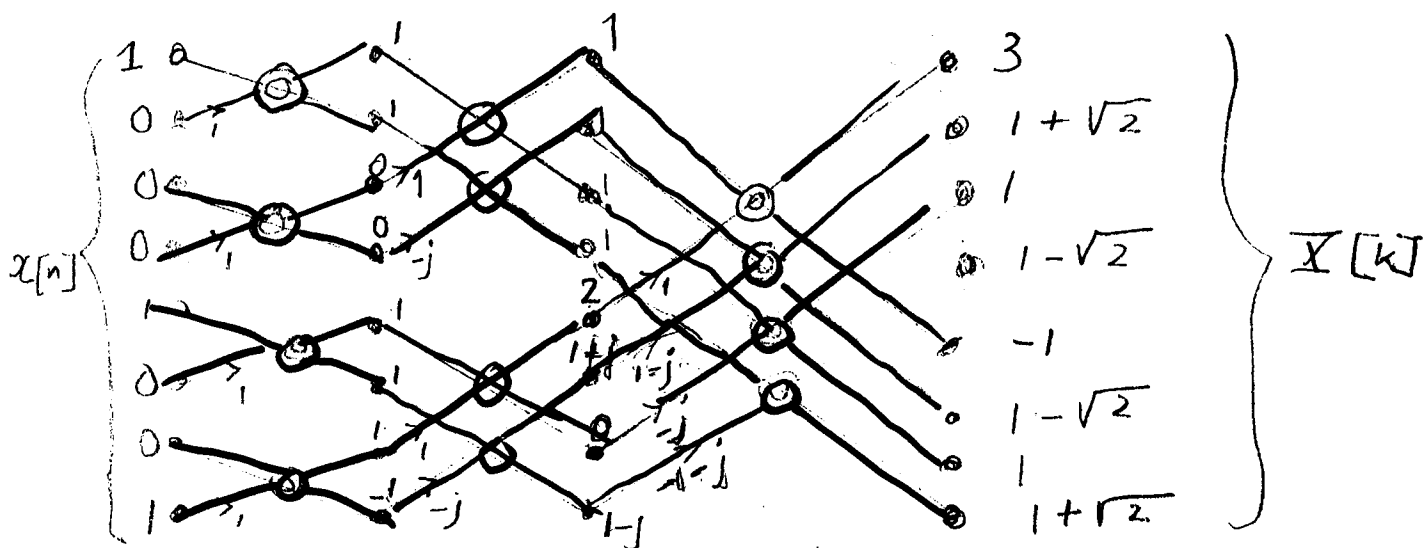
So at the start of the FFT algorithm reorders the data stream by bit reversal and performs the 2-Point DFT on the newly ordered streams.

This is clear from the soln of SP-1

Prob. SP-4

(a) Detailed Calculation of FFT

is depicted in the below signal flow graph



(Note that the input data stream is shuffled automatically by bit reversed or by the procedure explained in lecture)

(b)

$|X[k]|$



Prob. SP-2

8 (CONT)

(a)

$$x[n] = 0 \ 1 \ 1 \ 1$$

$$y[n] = 1 \ 2 \ 1$$

To perform circular convolution
extend above to strings by
zero padding to length

$$N = 4 + 3 - 1 = 6$$

Extended String

$$X[n] = 0 \ 1 \ 1 \ 1 \ 0 \ 0$$

$$y[n] = 1 \ 2 \ 1 \ 0 \ 0 \ 0$$

CONVOLUTION:

Using method in Lec. Notes #6,

pp. 18-20

Write

$$\hat{y} = 1 \ 0 \ 0 \ 0 \ 1 \ 2$$

then the circulant matrix

$$Y = \begin{pmatrix} 1 & 0 & 0 & 0 & 1 & 2 \\ 2 & 1 & 0 & 0 & 0 & 0 \\ 1 & 2 & 1 & 0 & 0 & 0 \\ 0 & 1 & 2 & 1 & 0 & 0 \\ 0 & 0 & 1 & 2 & 1 & 0 \\ 0 & 0 & 0 & 1 & 2 & 1 \end{pmatrix}$$

The convolution values are given by the vector z below

$$\begin{pmatrix} z[0] \\ z[1] \\ z[2] \\ z[3] \\ z[4] \\ z[5] \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 & 1 & 2 \\ 2 & 1 & 0 & 0 & 0 & 1 \\ 1 & 2 & 1 & 0 & 0 & 0 \\ 0 & 1 & 2 & 1 & 0 & 0 \\ 0 & 0 & 1 & 2 & 1 & 0 \\ 0 & 0 & 0 & 1 & 2 & 1 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \\ 1 \\ 1 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \\ 3 \\ 4 \\ 3 \\ 1 \end{pmatrix}$$

Using DFT Method

- First find DFTs of extended x and y by Matrix Method

$$\underline{X} = \begin{pmatrix} 1 & e^{-j\pi/3} & e^{-j2\pi/3} & 1 & e^{-j4\pi/3} & e^{-j\pi/3} \\ 1 & e^{-j2\pi/3} & e^{-j4\pi/3} & -1 & e^{-j\pi/3} & e^{-j4\pi/3} \\ 1 & -1 & e^{-j2\pi/3} & -1 & e^{j2\pi/3} & e^{-j\pi/3} \\ 1 & e^{j\pi/3} & e^{-j2\pi/3} & 1 & e^{j2\pi/3} & e^{-j\pi/3} \\ 1 & e^{j\pi/3} & e^{-j4\pi/3} & -1 & e^{-j2\pi/3} & e^{-j\pi/3} \end{pmatrix} \begin{pmatrix} 0 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \end{pmatrix}$$

$$= \begin{pmatrix} 3 & e^{-j\pi/3} + e^{-j2\pi/3} & -1 \\ e^{-j2\pi/3} + e^{-j4\pi/3} & +1 & \\ -1 & e^{j2\pi/3} + e^{-j2\pi/3} & +1 \\ e^{j\pi/3} + e^{-j4\pi/3} & -1 & \end{pmatrix} = \begin{pmatrix} 3 \\ -1-j \\ 0 \\ -1 \\ 0 \\ -1 \end{pmatrix}$$

$$Y = \begin{pmatrix} 1 & 1 & 1 & 1 & e^{-j\frac{4\pi}{3}} & e^{-j\frac{5\pi}{3}} \\ 1 & e^{-j\frac{\pi}{3}} & e^{-j\frac{2\pi}{3}} & -1 & e^{-j\frac{4\pi}{3}} & e^{-j\frac{5\pi}{3}} \\ 1 & e^{-j\frac{2\pi}{3}} & e^{-j\frac{4\pi}{3}} & 1 & e^{-j\frac{5\pi}{3}} & e^{-j\frac{4\pi}{3}} \\ 1 & -1 & 1 & -1 & 1 & -1 \\ 1 & e^{j\frac{\pi}{3}} & e^{j\frac{2\pi}{3}} & 1 & e^{j\frac{4\pi}{3}} & e^{j\frac{5\pi}{3}} \\ 1 & e^{j\frac{2\pi}{3}} & e^{j\frac{4\pi}{3}} & -1 & e^{-j\frac{5\pi}{3}} & e^{-j\frac{4\pi}{3}} \end{pmatrix} = \begin{pmatrix} 1 \\ 2 \\ 1 \\ 0 \\ 0 \\ 0 \end{pmatrix}$$

$$= \begin{pmatrix} 1 & & & & & \\ 1 + 2e^{-j\frac{\pi}{3}} + e^{-j\frac{2\pi}{3}} & & & & & \\ 1 + 2e^{-j\frac{2\pi}{3}} + e^{-j\frac{4\pi}{3}} & & & & & \\ 0 & & & & & \\ 1 + 2e^{j\frac{2\pi}{3}} + e^{j\frac{4\pi}{3}} & & & & & \\ 1 + 2e^{j\frac{\pi}{3}} + e^{j\frac{2\pi}{3}} & & & & & \end{pmatrix}$$

$$Z = \begin{pmatrix} (3)(1) & & & & & \\ -(1+j)(1 + 2e^{-j\frac{\pi}{3}} + e^{-j\frac{2\pi}{3}}) & & & & & \\ 0 & & & & & \\ -0 & & & & & \\ 0 & & & & & \\ -(1 + 2e^{j\frac{\pi}{3}} + e^{j\frac{2\pi}{3}}) & & & & & \end{pmatrix} = \begin{pmatrix} Z[6] \\ Z[0] \\ Z[2] \\ Z[3] \\ Z[4] \\ Z[5] \end{pmatrix}$$

HWK #8 Solns (CONT)

Invert z by

$$\begin{pmatrix} 0 \\ 1 \\ 3 \\ 4 \\ 3 \\ 1 \end{pmatrix} = z = \begin{pmatrix} z[0] \\ z[1] \\ \vdots \\ z[5] \end{pmatrix} = \frac{1}{6} \begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 \\ e^{-j\pi/3} & e^{j\pi/3} & e^{j2\pi/3} & e^{-j2\pi/3} & e^{j4\pi/3} & e^{-j4\pi/3} \\ e^{-j2\pi/3} & e^{j2\pi/3} & e^{j4\pi/3} & e^{-j4\pi/3} & e^{j6\pi/3} & e^{-j6\pi/3} \\ 1 & -1 & 1 & -1 & 1 & -1 \\ e^{-j\pi/3} & e^{j\pi/3} & e^{j2\pi/3} & e^{-j2\pi/3} & e^{j4\pi/3} & e^{-j4\pi/3} \\ e^{-j2\pi/3} & e^{j2\pi/3} & e^{j4\pi/3} & e^{-j4\pi/3} & e^{j6\pi/3} & e^{-j6\pi/3} \end{pmatrix} \begin{bmatrix} z \\ 1 \\ \vdots \\ 1 \\ z \end{bmatrix}$$

This is the convolution which will agree with those obtained by time domain procedure in p. 9 of these HWK #8 solution

$$\underline{W}^{-1} \underline{W}^*$$

(See eqns. (9) and (10) in Lec Notes #6)

I will ship (b) and move on to SR3.

Solns. to HWK #8 (CONT)

Prob. SP-3

$$(a) \quad x[n] = 0 \ 1 \ 1 \ 1$$

$$y[n] = 1 \ 2 \ 1$$

Extended Strings

$$x[n] = 0 \ 1 \ 1 \ 1 \ 0 \ 0$$

$$y[n] = 1 \ 2 \ 1 \ 0 \ 0 \ 0$$

CORRELATION

Circulant Matrix $\tilde{Y}$ (See p. 25 of Lec Notes #6)

$$\tilde{Y} = \begin{pmatrix} 1 & 2 & 1 & 0 & 0 & 0 \\ 2 & 1 & 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 1 & 2 \\ 0 & 0 & 0 & 1 & 2 & 1 \\ 0 & 0 & 1 & 2 & 1 & 0 \\ 0 & 1 & 2 & 1 & 0 & 0 \end{pmatrix}$$

$$\text{Correlation Vector} \begin{pmatrix} z[0] \\ z[1] \\ z[2] \\ z[3] \\ z[4] \\ z[5] \end{pmatrix} = \begin{pmatrix} 1 & 2 & 1 & 0 & 0 & 0 \\ 2 & 1 & 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 1 & 2 \\ 0 & 0 & 0 & 1 & 2 & 1 \\ 0 & 0 & 1 & 2 & 1 & 0 \\ 0 & 1 & 2 & 1 & 0 & 0 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \\ 1 \\ 1 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} 3 \\ 1 \\ 0 \\ 1 \\ 3 \\ 4 \end{pmatrix}$$

Using DFT Method

I will outline the procedure

→ Obtain the same DFT's as in SP-2

$$X = W x = \begin{pmatrix} X[0] \\ X[1] \\ \vdots \\ X[5] \end{pmatrix}$$

$$Y = W y = \begin{pmatrix} Y[0] \\ Y[1] \\ \vdots \\ Y[5] \end{pmatrix}$$

→ Obtain the DFT of Correlation vector*

$$Z = \begin{pmatrix} z[0] \\ z[1] \\ \vdots \\ z[5] \end{pmatrix} = \begin{pmatrix} X[0] Y[0] \\ X^*[1] Y[1] \\ \vdots \\ X^*[5] Y[5] \end{pmatrix} = \begin{pmatrix} z[0] \\ z[1] \\ \vdots \\ z[5] \end{pmatrix}$$

Homework #8 (CONT)

Prob SP-4

$$h[n] = [1, 2, 1]$$

$$x[n] = [x_0, x_1, \dots, x_{95}]$$

In order to perform FFTs
You need block length after 0-padding
equal to a power of 2

Since $h[n]$ has length $3 = M$
you will need $x[n]$ blocks to
have length 2 to achieve
 $N_1 + N_2 - 1 = 3 + 2 - 1 = 4$

So sensible choices are

$$N_2 = 2 \rightarrow \text{block length } 4$$

$$N_2 = 6 \rightarrow \text{block length } 8$$

Total no. of calculations
1st choice $N_2 = 2$ block length $N = 4$ (24 blocks)

$$\text{Calc. for each block} = \frac{N}{2} \log_2 N = \frac{4}{2} \log_2 4 = (2)(2) = 4$$

$$\text{Total no. of calcs} = (24)(4) = \boxed{96} \text{ (12 blocks)}$$

2nd. choice $N_2 = 6$ block length = 8

$$\text{Calc/block} = \frac{8}{2} \log_2 8 = (4)(3) = 12$$

$$\text{Total no. of calcs} = (12)(8) = \boxed{96}$$

Hunk Silm # 8 (cont)

Both block sizes lead to the same amount of calculations

We choose block size = 4

Realization / Architecture

