

Department of Systems Engineering
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**SYST 302: Systems Methodology
and Design II #5**

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Outline

- **Resource Allocation**
- **Optimization Theory**
- **Constraint and Unconstraint Optimization**
- **Numerical Optimization Techniques**

Resource Allocation and Optimization

- **Resource Allocation:** distribute resources to users so that the highest level of a specified objective can be achieved
 - Factory organization for maximum productivity
 - Spacecraft flight plan to minimize fuel consumption
 - Security and safety procedures to minimize risks and hazards
- **Optimization Problems:** to achieve a performance in the extreme within the resource available
 - Given budget resource, optimize performance (maximize reliability, minimize risk)
 - Given specification of performance, minimize the capital required

Optimization

- **Problem Formulation:** translate the desired system performance specifications and constraints into a mathematical form
 - Set up objective function (revenue, yield, cost, reliability, etc.)
 - Identify control variables
 - Specify constraints
- **Optimization Techniques:**
 - Unconstrained optimization
 - Constrained optimization
 - Linear/nonlinear programming
 - Dynamic programming

Unconstrained Optimization Single Variables

Problem: Minimize $f(x)$,

second - order Taylor series expansion

$$f(x) = f(x_0) + f_x(x_0)(x - x_0) + \frac{1}{2} f_{xx}(x_0)(x - x_0)^2 + E(x - x_0)$$

where

$$f_x(x_0) \equiv \frac{\partial f(x_0)}{\partial x} \text{ is the first order derivative}$$

$$f_{xx}(x_0) = \frac{\partial^2 f(x_0)}{\partial x^2} \text{ is the second order derivative}$$

$$\text{Ex: } f(x) = 2x^3 - 3x^2 + 2$$

$$f_x(x) = 6x^2 - 6x, \quad f_{xx}(x) = 12x - 6$$

Unconstrained Optimization

Multiple variables

Problem: Minimize $f(\mathbf{x})$, where $\mathbf{x} = [x_1, \dots, x_m]^T$

second - order Taylor series expansion

$$f(\mathbf{x}) = f(\mathbf{x}_0) + f_{\mathbf{x}}(\mathbf{x}_0)(\mathbf{x} - \mathbf{x}_0) + \frac{1}{2}(\mathbf{x} - \mathbf{x}_0)^T f_{\mathbf{xx}}(\mathbf{x}_0)(\mathbf{x} - \mathbf{x}_0) + E(\mathbf{x} - \mathbf{x}_0)$$

where

$$f_{\mathbf{x}}(\mathbf{x}_0) \equiv \nabla f(\mathbf{x}_0) = [f_{x_1}(\mathbf{x}_0), f_{x_2}(\mathbf{x}_0), \dots, f_{x_m}(\mathbf{x}_0)]$$

$$= \left[\frac{\partial f}{\partial x_1}, \frac{\partial f}{\partial x_2}, \dots, \frac{\partial f}{\partial x_m} \right]_{\mathbf{x}=\mathbf{x}_0} \text{ is called } \textit{gradient}$$

$$f_{\mathbf{xx}}(\mathbf{x}_0) = \nabla^2 f(\mathbf{x}_0) = \begin{bmatrix} f_{x_1 x_1} & \cdots & f_{x_1 x_m} \\ \vdots & \ddots & \vdots \\ f_{x_m x_1} & \cdots & f_{x_m x_m} \end{bmatrix} = \left[\frac{\partial^2 f}{\partial x_i \partial x_j} \right]_{\mathbf{x}=\mathbf{x}_0} \text{ is called } \textit{Hessian}$$

Example

$$f(x_1, x_2) = x_1^2 - 6x_1 + 2x_2^2 - 8x_2 + 30$$

$$f_{\mathbf{x}}(\mathbf{x}_0) \equiv \nabla f(\mathbf{x}_0) = \left[\frac{\partial f}{\partial x_1}, \frac{\partial f}{\partial x_2} \right]_{\mathbf{x}=\mathbf{x}_0} = [2x_1 - 6 \quad 4x_2 - 8]$$

$$f_{\mathbf{xx}}(\mathbf{x}_0) = \nabla^2 f(\mathbf{x}_0) = \left[\begin{array}{cc} \frac{\partial^2 f}{\partial x_1^2} & \frac{\partial^2 f}{\partial x_1 \partial x_2} \\ \frac{\partial^2 f}{\partial x_1 \partial x_2} & \frac{\partial^2 f}{\partial x_2^2} \end{array} \right]_{\mathbf{x}=\mathbf{x}_0} = \begin{bmatrix} 2 & 0 \\ 0 & 4 \end{bmatrix}$$

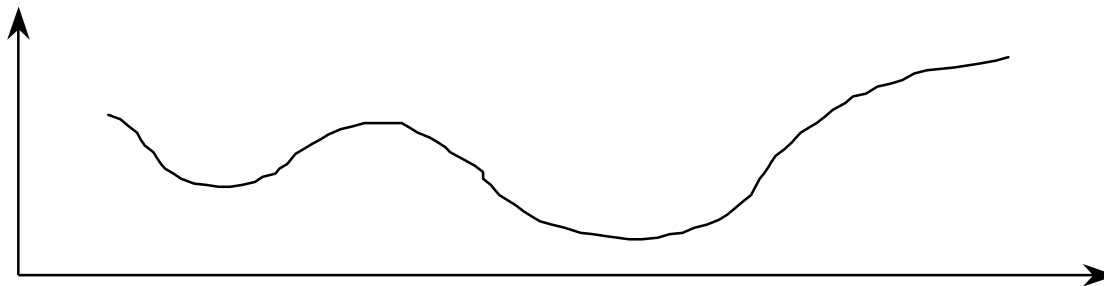
Local and Global Minimum

Local Minimum:

A point x_0 is said to be a relative or local minimum if there exists an $\varepsilon > 0$, such that $f(x) \geq f(x_0)$ for all x , such that $|x - x_0| < \varepsilon$ and $x \in \Omega$. This definition says that if we can draw an arbitrarily small box around our point u_0 and no other values of $f(x)$ within that small box have a lower value than $f(x_0)$, then we have a local minimum.

Global Minimum:

A point x^* is said to be a global minimum if $f(x) \geq f(x^*)$ for all $x \in \Omega$.



Necessary and Sufficient Conditions Single Variable

Local Minimum: (Maximun)

Necessary consitions

first - order $f_x(x_0) \equiv \frac{\partial f(x_0)}{\partial x} = 0$

second - order $f_{xx}(x_0) = \frac{\partial^2 f(x_0)}{\partial x^2} \geq 0 \ (\leq 0)$

Sufficient consitions:

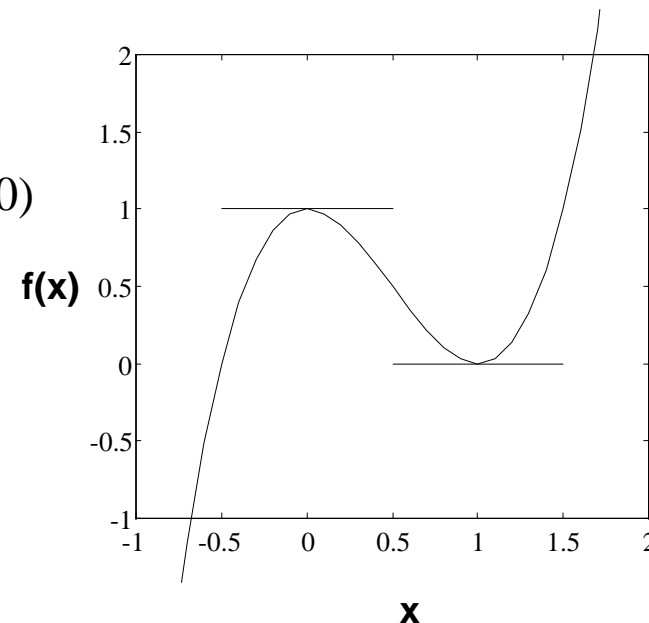
first - order $f_x(x_0) \equiv \frac{\partial f(x_0)}{\partial x} = 0$

second - order $f_{xx}(x_0) = \frac{\partial^2 f(x_0)}{\partial x^2} > 0 \ (< 0)$

Ex: $f(x) = 2x^3 - 3x^2 + 2$

$f_x(x) = 6x^2 - 6x = 0 \Rightarrow x = 1 \text{ or } 0$

$f_{xx}(x) = 12x - 6 = 6 \text{ or } -6$



Necessary and Sufficient Conditions Multiple Variables

Local Minimum: (Maximum)

Necessary conditions

first - order $\nabla f(\mathbf{x}_0) = \mathbf{0}$ (\mathbf{x}_0 : critical point)

second - order $f_{\mathbf{xx}}(\mathbf{x}_0) \equiv Q$ is positive semidefinite ($\mathbf{x}^T Q \mathbf{x} \geq 0, \forall \mathbf{x} \neq 0$)

Sufficient conditions:

first - order $\nabla f(\mathbf{x}_0) = \mathbf{0}$

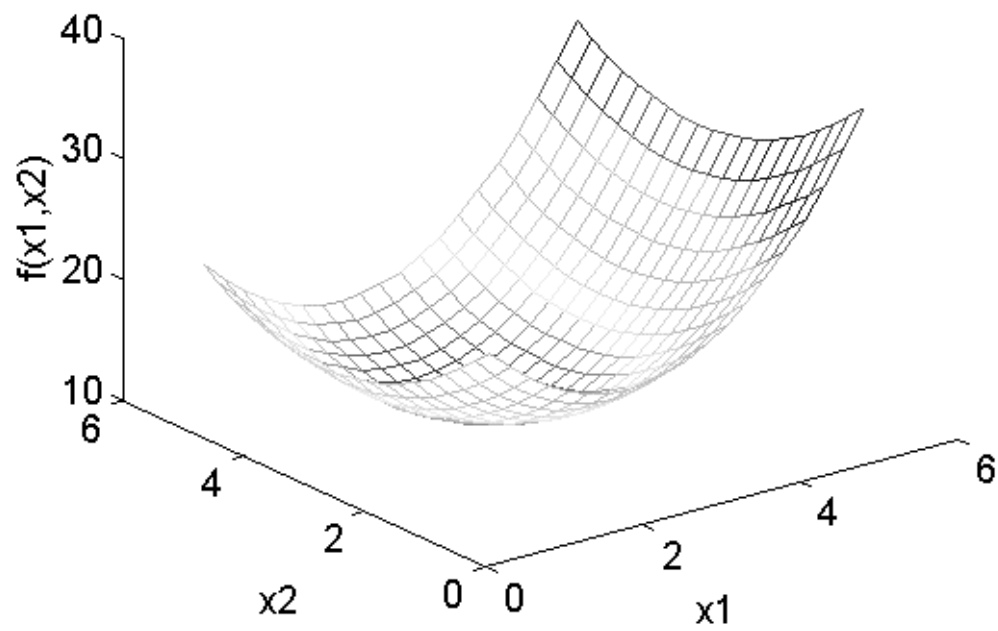
second - order $f_{\mathbf{xx}}(\mathbf{x}_0) \equiv Q$ is positive definite ($\mathbf{x}^T Q \mathbf{x} > 0, \forall \mathbf{x} \neq 0$)

Ex: $f(x_1, x_2) = x_1^2 - 6x_1 + 2x_2^2 - 8x_2 + 30$

$f_{\mathbf{x}}(\mathbf{x}_0) \equiv \nabla f(\mathbf{x}_0) = [2x_1 - 6 \quad 4x_2 - 8] = [0 \ 0] \Rightarrow x_1 = 3, x_2 = 2$

$f_{\mathbf{xx}}(\mathbf{x}_0) = \nabla^2 f(\mathbf{x}_0) = Q = \begin{bmatrix} 2 & 0 \\ 0 & 4 \end{bmatrix} \Rightarrow Q \text{ is positive definite} \Rightarrow \text{minimum}$

Example



Bridge Design Example

Total cost = cost of superstructure + cost of piers

$$C_T = (AS + B)LC_S + \left(\frac{L}{S} + 1\right)C_p$$

$$\frac{\partial C_T}{\partial S} = AL C_S - \frac{LC_p}{S^2} = 0 \Rightarrow S^* = \sqrt{\frac{C_p}{AC_S}}$$

Note: Need to choose $n=L/S^*$ to a closest integer

L : bridge length (feet)

W : superstructure weight (pounds per foot, $W = AS + B$)

S : span between piers (foot)

C_S : cost of superstructure (\$ per pound)

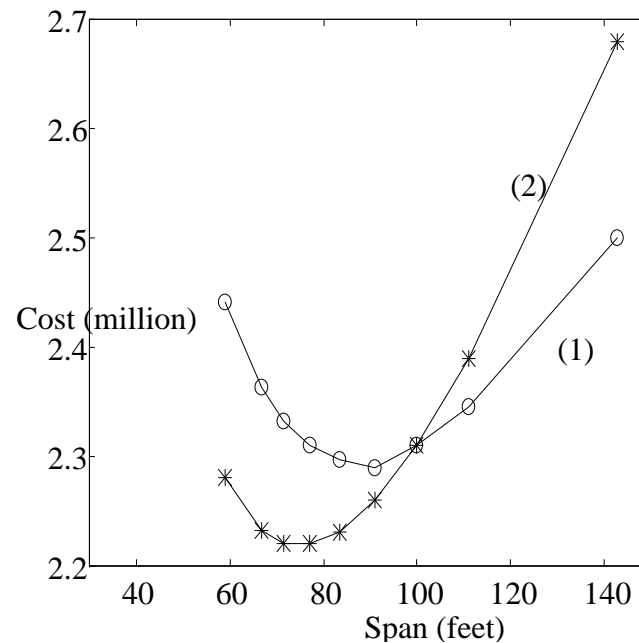
C_p : cost of piers (\$ per pier)

Alternative Design

Given $L = 1,000$ $C_s = 0.65$, $C_p = 80k$

$$(1) W_1 = 16S + 600 S^* = 87.7\text{ft}$$

$$(2) W_2 = 22S \quad S^* = 74.8\text{ft}$$



Conductor Design Example

cost model

Total cost = Power loss cost + Investment cost

$C_T = f$ (cross-sectional area, electricity cost,
unit cost of the conductor,...)

Issue: cross section \uparrow , installed cost \uparrow , power loss cost \downarrow

$$\text{Power loss cost /year} = C_e I^2 R \frac{Hrs}{1000A}$$

$$\text{Investment Cost (annual)} = C_T = C_i \left(\frac{A}{P} \right) + (C_m - F_m) W_m \left(\frac{A}{P} \right) + F_m W_m i$$

where $W_m = LAD_m / 144$, A : area of the cross section (in²)

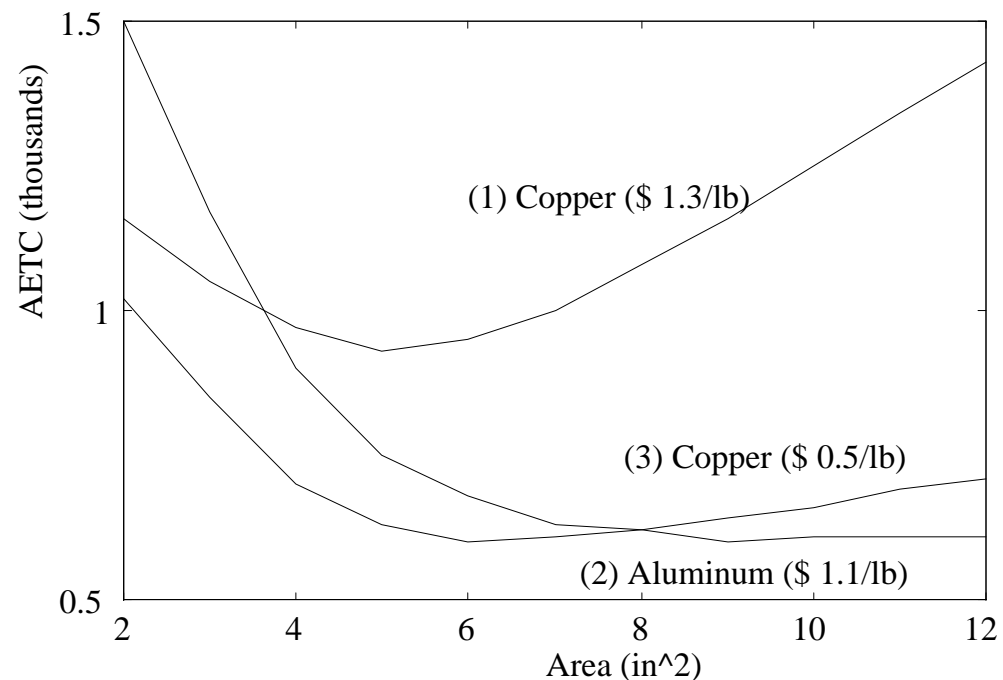
C_i : installation cost, D_m : material density (weight per cubic foot)

C_m, F_m : present and future value of the material per unit weight

Cost Curves

$$\frac{\partial C_T}{\partial A} = 0 \Rightarrow A^* = \sqrt{\frac{C_e I^2 R H / 1000}{(C_m - F_m) \left(\frac{A}{P}\right) L \frac{D_m}{144} + i F_m L \frac{D_m}{144}}$$

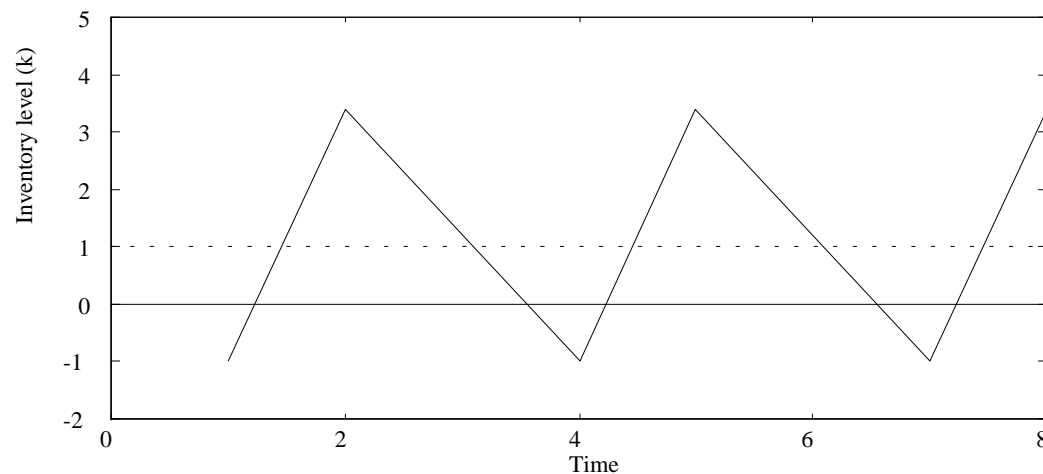
$$\frac{\partial^2 C_T}{\partial A^2} = 2C_e I^2 R \frac{H}{1000A^3} > 0 \Rightarrow \min$$



Inventory Decision Example

Total cost = item cost + procurement cost + holding cost
+ shortage cost

$C_T = f(\text{procurement level, procurement quantity, demand rate, lead time, replenishment rate, ...})$



Assumptions: replenishment rate is greater than demand rate and that unsatisfied demand is not lost

Inventory Decision Cost Model

C_i : item unit cost C_p : procurement cost per procurement

D : demand rate in units per period

L : procurement level Q : procurement quantity

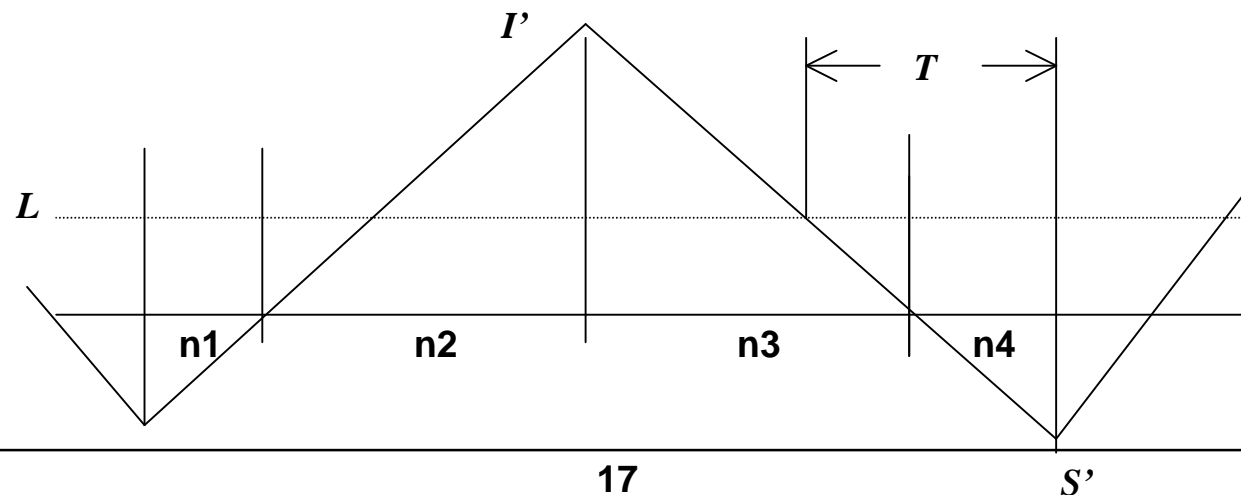
$N = Q / D = n_1 + n_2 + n_3 + n_4$: number of periods per cycle

$I = (n_2 + n_3)I' / 2$: total inventory per cycle

$S = (n_4 + n_1)S' / 2$: total shortage per cycle

C_h : unit holding cost per period

C_s : unit shortage cost per period



Inventory Decision Cost Model

item cost $= C_i D$

procurement cost $= C_p / N = C_p \frac{D}{Q}$

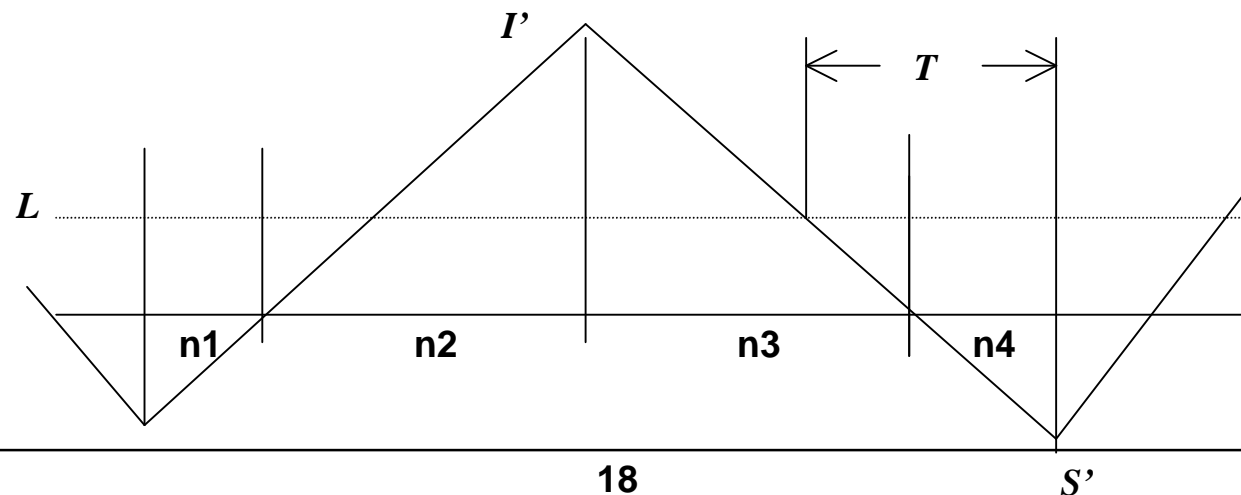
holding cost $= C_h I / N = C_h \frac{D}{Q} \frac{(n_2 + n_3)I'}{2}$

shortage cost $= C_s S / N = C_s \frac{D}{Q} \frac{(n_1 + n_4)S'}{2}$

where $(n_1 + n_2)(R - D) = (n_3 + n_4)D$; $n_1 + n_2 = Q / R$

$$n_3 + n_4 = \frac{I' - L}{D} + T; \quad S' = DT - L$$

Total cost $= C_T = C_i D + C_p / N + C_h I / N + C_s S / N$



Optimal Inventory Decision

$$\text{solve } \frac{\partial C_T}{\partial L} = 0, \frac{\partial C_T}{\partial Q} = 0$$

$$\Rightarrow Q^* = \sqrt{\frac{1}{1 - D/R} \left(\frac{2C_p D}{C_h} + \frac{2C_p D}{C_s} \right)}$$

$$L^* = DT - \sqrt{\left(1 - \frac{D}{R}\right) \frac{2C_p D}{C_s(1 + C_s/C_h)}}$$

$$C_T^* = C_i D + \sqrt{\left(1 - \frac{D}{R}\right) \frac{2C_p C_h C_s D}{C_s + C_h}}$$

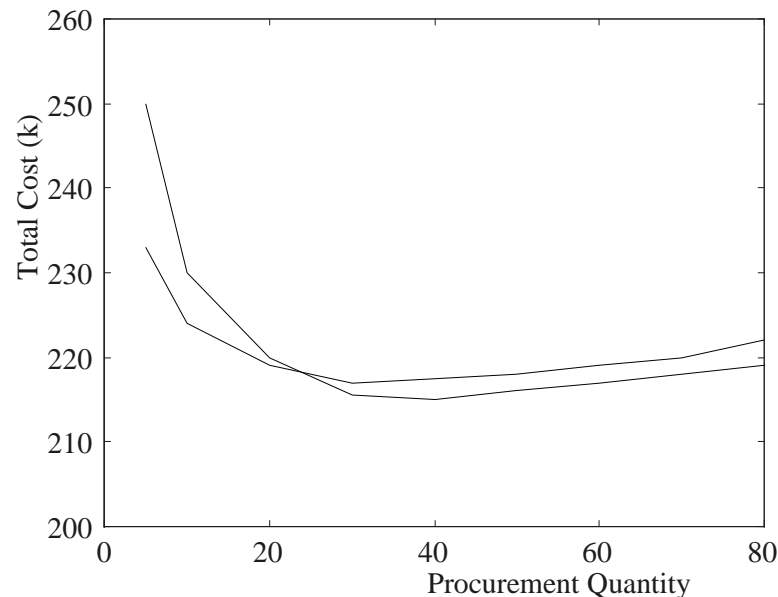
sufficient condition: $\begin{bmatrix} \frac{\partial^2 C_T}{\partial Q^2} & \frac{\partial^2 C_T}{\partial Q \partial L} \\ \frac{\partial^2 C_T}{\partial Q \partial L} & \frac{\partial^2 C_T}{\partial L^2} \end{bmatrix}$ is positive definite

Space Constrained Operations

Warehouse restriction : max inventory is constrained

Ex : unconstrained $\Rightarrow Q^* = 43.04, L^* = 20.48$

constrained (I') $\Rightarrow Q^* = 30.57, L^* = 15.6$



Numerical Optimization Techniques

- **Motivation:** analytical solution may not be feasible
 - Cost function may not be differentiable
 - Evaluating the expression for the derivative may be difficult
- **Algorithm Properties:**
 - *Iterative* in nature
 - *Decent* in sequence
 - *Global convergent* is desirable but often not achievable
- **Numerical Techniques:**
 - Search methods
 - Gradient methods

Search Methods

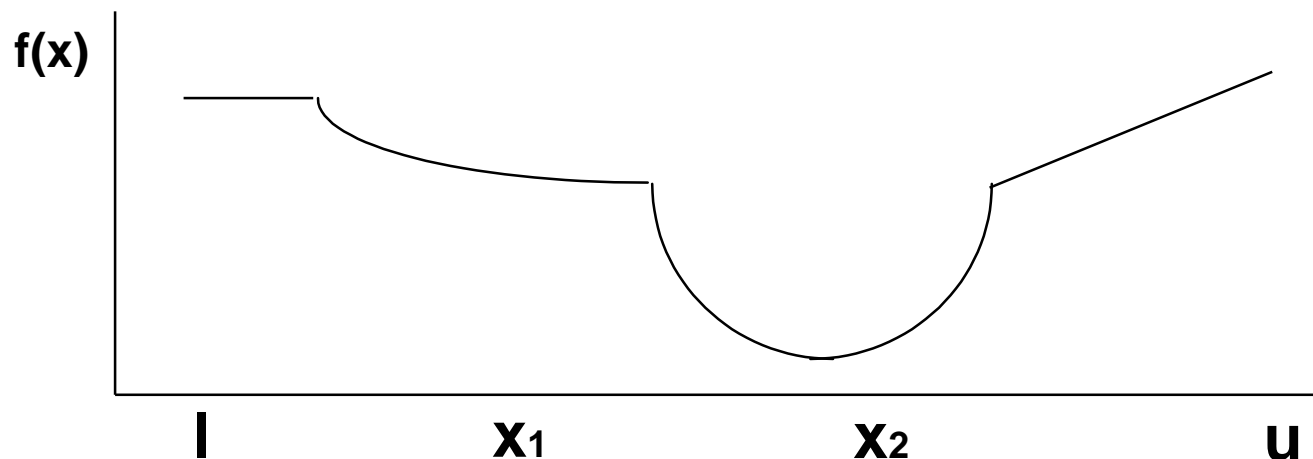
Constructing shrinking bounding intervals

Fibonacci , Golden section

Often improved by curve - fitting line search

Basic property of unimodal function :

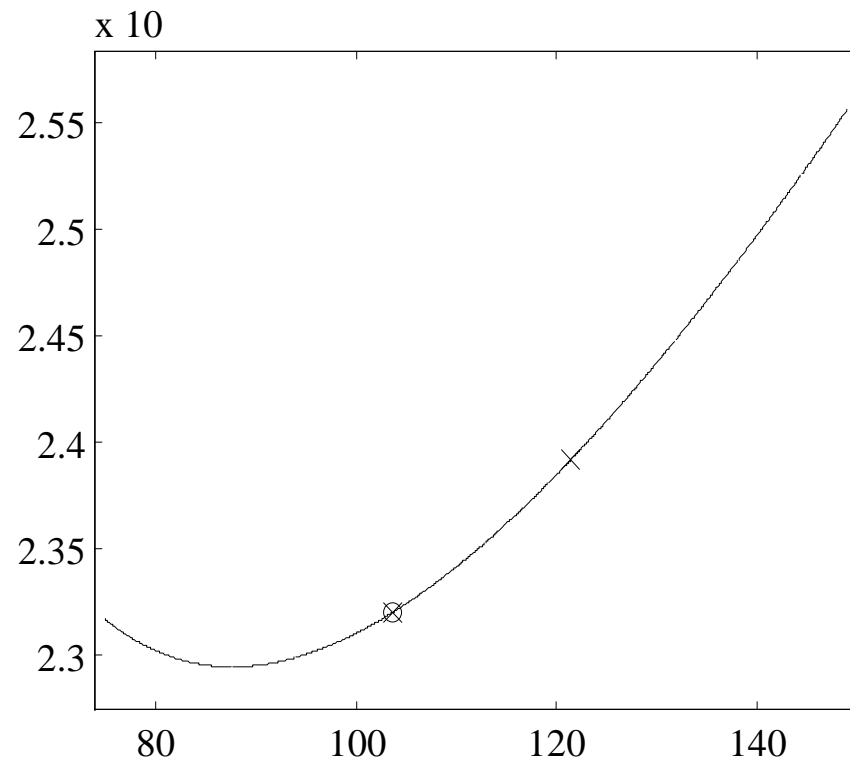
- (1) if $f(x_1) > f(x_2) \Rightarrow x^* \in [x_1, u]$
- (2) if $f(x_1) < f(x_2) \Rightarrow x^* \in [l, x_2]$
- (3) if $f(x_1) = f(x_2) \Rightarrow x^* \in [x_1, x_2]$



Example - Golden Section Search

Bridge design example

$$f(s) = 10400s + 80000000 / s + 470000$$



Gradient Methods

- **Steepest Decent:** oldest and most widely known methods for minimizing multi-variable functions
 - Simple and not difficult to analyze
 - Many modification on the basic technique exist for better convergence properties
- **The Method:**

To minimize $f(\mathbf{x})$

$\mathbf{x}_{k+1} = \mathbf{x}_k - \alpha_k \mathbf{g}_k$ where $\mathbf{g}_k = \nabla f(\mathbf{x}_k)$

and α_k is a nonnegative scalar minimizing $f(\mathbf{x}_k - \alpha_k \mathbf{g}_k)$

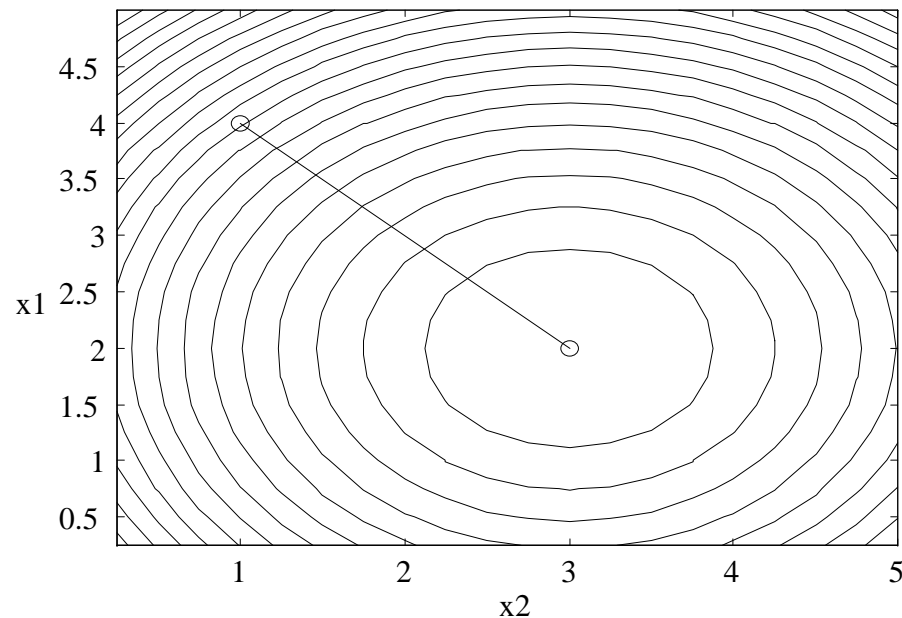
Example I

$$\text{Ex: } f(x_1, x_2) = x_1^2 - 4x_1 + x_2^2 - 6x_2 + 30; \quad x_1(0) = 4, x_2(0) = 1$$

$$\nabla f(\mathbf{x}(0)) = \begin{bmatrix} 2x_1 - 4 & 2x_2 - 6 \end{bmatrix}_{\mathbf{x}(0)} = \begin{bmatrix} 4 & -4 \end{bmatrix}, \quad x_1(1) = x_1(0) - 4\alpha, x_2(1) = x_2(0) + 4\alpha$$

$$f(\mathbf{x}(1)) = (4 - 4\alpha)^2 - 4(4 - 4\alpha) + (1 + 4\alpha)^2 - 6(1 + 4\alpha) + 30 = 32\alpha^2 - 32\alpha + 25 \equiv h(\alpha)$$

$$\min h(\alpha) \Rightarrow \alpha = 0.5 \Rightarrow x_1(1) = 2, x_2(1) = 3$$



Example II

$$Ex : f(x_1, x_2) = 2x_1^2 - 6x_1 + x_2^2 - 2x_1x_2; x_1(0) = 0, x_2(0) = 0$$

$$\nabla f(\mathbf{x}(0)) = \begin{bmatrix} 4x_1 - 6 - 2x_2 & 2x_2 - 2x_1 \end{bmatrix}_{\mathbf{x}(0)} = \begin{bmatrix} -6 & 0 \end{bmatrix}, x_1(1) = x_1(0) + 6\alpha, x_2(1) = x_2(0) + 0\alpha$$

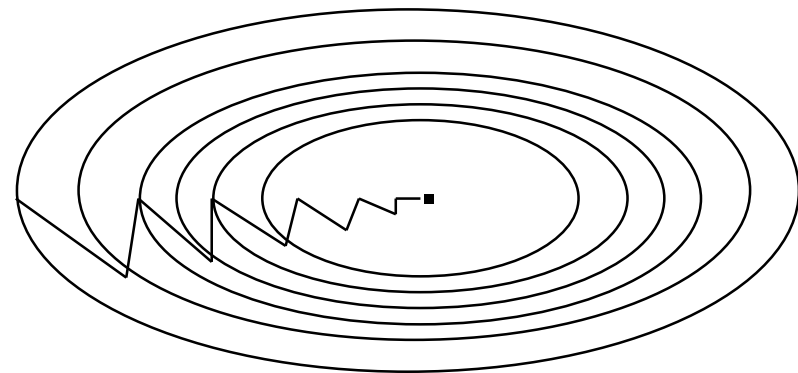
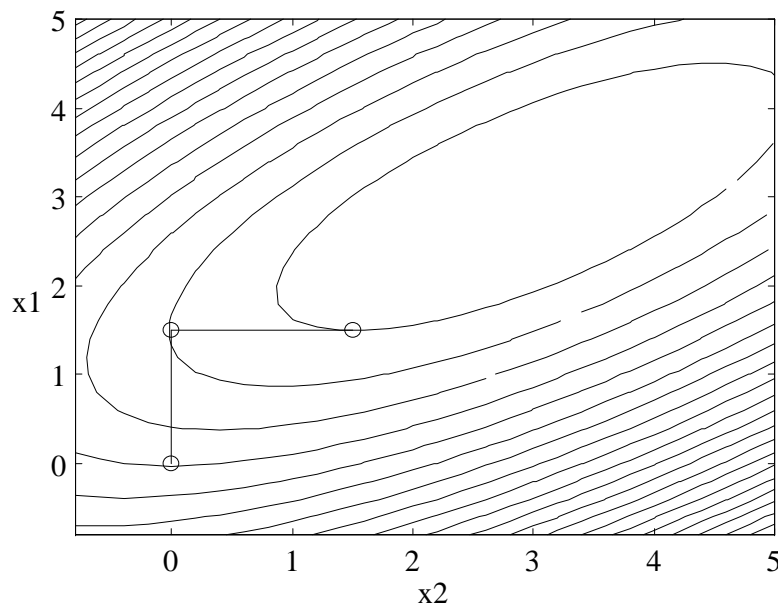
$$f(\mathbf{x}(1)) = 2(6\alpha)^2 - 6(6\alpha) + (0)^2 - 2(6\alpha)(0) = 72\alpha^2 - 36\alpha \equiv h(\alpha)$$

$$\min h(\alpha) \Rightarrow \alpha = 0.25 \Rightarrow x_1(1) = 1.5, x_2(1) = 0$$

$$\nabla f(\mathbf{x}(1)) = \begin{bmatrix} 4x_1 - 6 - 2x_2 & 2x_2 - 2x_1 \end{bmatrix}_{\mathbf{x}(1)} = \begin{bmatrix} 0 & -3 \end{bmatrix}, x_1(2) = x_1(1) + 0\alpha, x_2(2) = x_2(1) + 3\alpha$$

$$f(\mathbf{x}(2)) = 2(1.5)^2 - 6(1.5) + (3\alpha)^2 - 2(1.5)(3\alpha) = 9\alpha^2 - 9\alpha - 4.5 \equiv h(\alpha)$$

$$\min h(\alpha) \Rightarrow \alpha = 0.5 \Rightarrow x_1(2) = 1.5, x_2(2) = 1.5$$



Newton's Method

Approximate the function locally by a quadratic function

$$f(\mathbf{x}) \cong f(\mathbf{x}_k) + \nabla f(\mathbf{x}_k)(\mathbf{x} - \mathbf{x}_k) + \frac{1}{2}(\mathbf{x} - \mathbf{x}_k)'[\nabla^2 f(\mathbf{x}_k)](\mathbf{x} - \mathbf{x}_k)$$

the right - hand side is minimized at

$$\mathbf{x}_{k+1} = \mathbf{x}_k - [\nabla^2 f(\mathbf{x}_k)]^{-1} \nabla f(\mathbf{x}_k) \Leftarrow \text{Newton's Method}$$

For real quadratic functions $f(\mathbf{x}) = \frac{1}{2} \mathbf{x}' Q \mathbf{x} - \mathbf{x}' \mathbf{b}$

$$\Rightarrow \mathbf{x}_{k+1} = \mathbf{x}_k - Q^{-1}(Q\mathbf{x}_k - \mathbf{b}) = Q^{-1}\mathbf{b}$$

reach solution in **one iteration**

Example

$$\text{Ex: } f(x_1, x_2) = 2x_1^2 - 6x_1 + x_2^2 - 2x_1x_2; x_1(0) = 0, x_2(0) = 0$$

$$\nabla f(\mathbf{x}(0)) = [4x_1 - 6 - 2x_2 \quad 2x_2 - 2x_1]_{\mathbf{x}(0)} = [-6 \quad 0], \nabla^2 f(\mathbf{x}(0)) = \begin{bmatrix} 4 & -2 \\ -2 & 2 \end{bmatrix}$$

$$[\nabla^2 f(\mathbf{x}(0))]^{-1} \nabla f(\mathbf{x}(0))^T = \begin{bmatrix} 0.5 & 0.5 \\ 0.5 & 1 \end{bmatrix} \begin{bmatrix} -6 \\ 0 \end{bmatrix} = \begin{bmatrix} -3 \\ -3 \end{bmatrix}$$

$$x_1(1) = x_1(0) - (-3) = 3, x_2(1) = x_2(0) - (-3) = 3$$

